OSQP.jl

A Julia wrapper for the Operator Splitting QP solver

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joint work with Goran Banjac,

Nicholas Moehle, Paul Goulart, Alberto Bemporad, Stephen Boyd

JuMP Developers Meetup, 13 Jun 2017

Why quadratic programming?

AN ALGORITHM FOR QUADRATIC PROGRAMMING

Marguerite Frank and Philip Wolfe[1]
Princeton University

A finite iteration method for calculating the solution of quadratic programming problems is described. Extensions to more general non-linear problems are suggested.

1. INTRODUCTION

The problem of maximizing a concave quadratic function whose variables are subject to linear inequality constraints has been the subject of several recent studies, from both the computational side and the theoretical (see Bibliography). Our aim here has been to develop a method for solving this non-linear programming problem which should be particularly well adapted to high-speed machine computation.

March 1956!





145
140
135
130
125
123.74
Chg
Tues.
142.78

First-order methods

Pros

- Warm starting
- Handle large-scale problems
- Embeddable

Cons

- Low accuracy solutions
- Don’t detect infeasibility
- Problem data dependent

General Purpose QP Solver

Based on first-order methods

Robust

Accurate

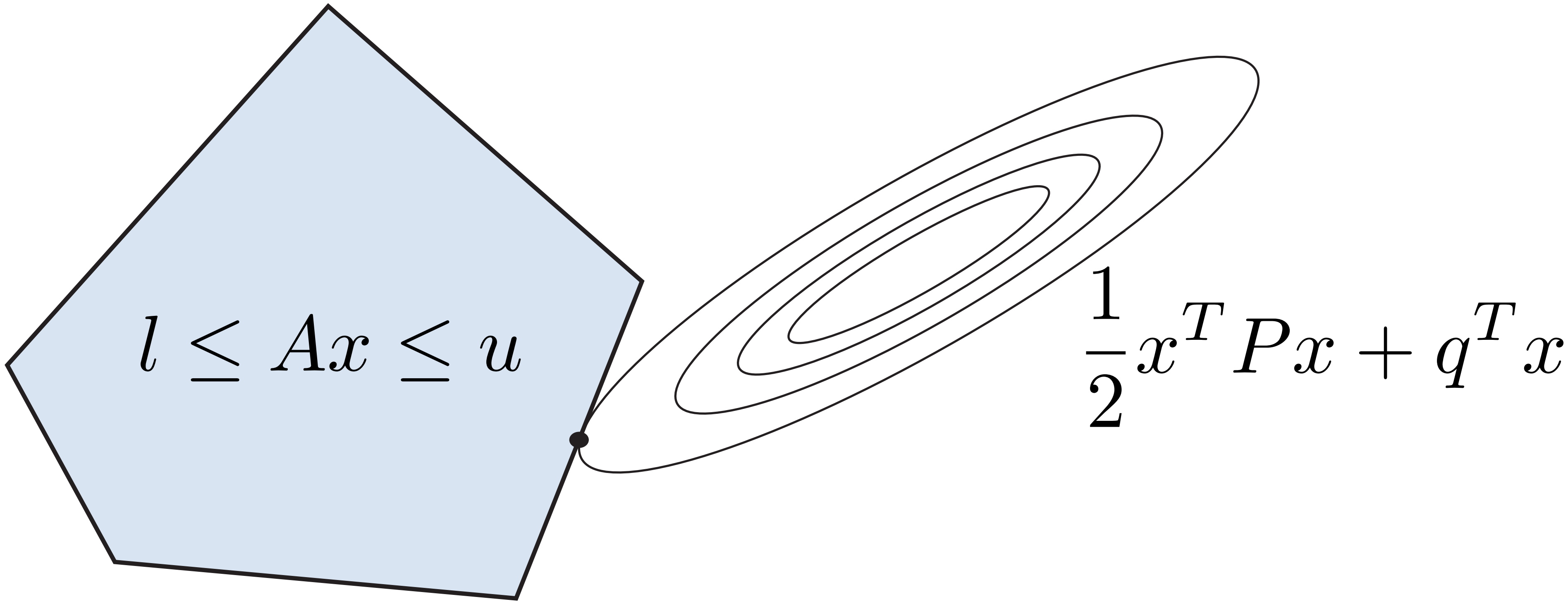
Detects Infeasibility

The OSQP Solver

The problem

Quadratic Program

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & l \leq Ax \leq u \end{array}$$



ADMM

minimize $f(x) + g(x)$ $\longrightarrow$ minimize $f(\tilde{x}) + g(x)$
subject to $\tilde{x} = x$

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1 $$\tilde{x}^{k+1} \leftarrow \operatorname*{argmin}_{\tilde{x}} \left(f(\tilde{x}) + \frac{\rho}{2} \left\| \tilde{x} - x^k + \frac{y^k}{\rho} \right\|^2 \right)$$

2 $$x^{k+1} \leftarrow \operatorname*{argmin}_{x} \left(g(x) + \frac{\rho}{2} \left\| x - \tilde{x}^{k+1} - \frac{y^k}{\rho} \right\|^2 \right)$$

3 $$y^{k+1} \leftarrow y^k + \rho \left(\tilde{x}^{k+1} - x^{k+1} \right)$$

How to split the QP?

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & Ax = z \\ & l \leq z \leq u \end{array}$$

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}\tilde{x}^T P \tilde{x} + q^T \tilde{x} + \mathcal{I}_{Ax=z}(\tilde{x}, \tilde{z}) + \mathcal{I}_{[l,u]}(z) \\ \text{subject to} & (\tilde{x}, \tilde{z}) = (x, z) \end{array}$$

How to split the QP?

$$\begin{array}{ll}\text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & Ax = z \\ & l \leq z \leq u\end{array} \qquad f$$

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How to split the QP?

minimize $\frac{1}{2}x^TPx + q^Tx$ — f

subject to $Ax = z$

$l \leq z \leq u$ — g

$$\underbrace{\frac{1}{2}\tilde{x}^TP\tilde{x} + q^T\tilde{x} + \mathcal{I}_{Ax=z}(\tilde{x},\tilde{z})}_{f} + \underbrace{\mathcal{I}_{[l,u]}(z)}_{g}$$

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subject to $(\tilde{x},\tilde{z}) = (x, z)$

OSQP

Problem

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & l \leq Ax \leq u \end{array}$$

Algorithm

1 $\begin{cases} (x^{k+1}, \nu^{k+1}) \leftarrow \text{solve} \begin{bmatrix} P + \sigma I & A^T \\ A & -\frac{1}{\rho} I \end{bmatrix} \begin{bmatrix} x^{k+1} \\ \nu^{k+1} \end{bmatrix} = \begin{bmatrix} \sigma x^k - q \\ z^k - \frac{1}{\rho} y^k \end{bmatrix} \\ \tilde{z}^{k+1} \leftarrow z^k + \frac{1}{\rho}(\nu^{k+1} - y^k) \end{cases}$

2 $\{ \; z^{k+1} \leftarrow \Pi\left(\tilde{z}^{k+1} + \frac{1}{\rho} y^k\right)$

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Algorithm

Linear system solve

$$(x^{k+1}, \nu^{k+1}) \leftarrow \text{solve} \begin{bmatrix} P + \sigma I & A^T \\ A & -\frac{1}{\rho} I \end{bmatrix} \begin{bmatrix} x^{k+1} \\ \nu^{k+1} \end{bmatrix} = \begin{bmatrix} \sigma x^k - q \\ z^k - \frac{1}{\rho} y^k \end{bmatrix}$$

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Easy operations

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minimize $\frac{1}{2}x^T P x + q^T x$
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Factorization caching

OSQP

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Warm starting

Factorization caching

OSQP

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$$\tilde{z}^{k+1} \leftarrow z^k + \frac{1}{\rho}(\nu^{k+1} - y^k)$$

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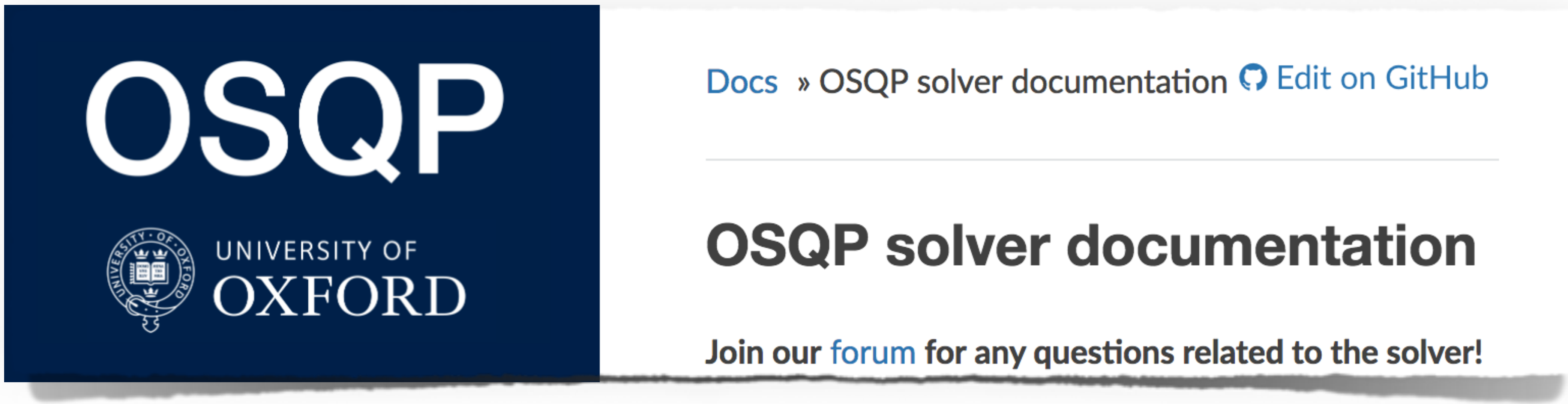
Solution polishing

Warm starting

Factorization caching

OSQP

`osqp.readthedocs.io`



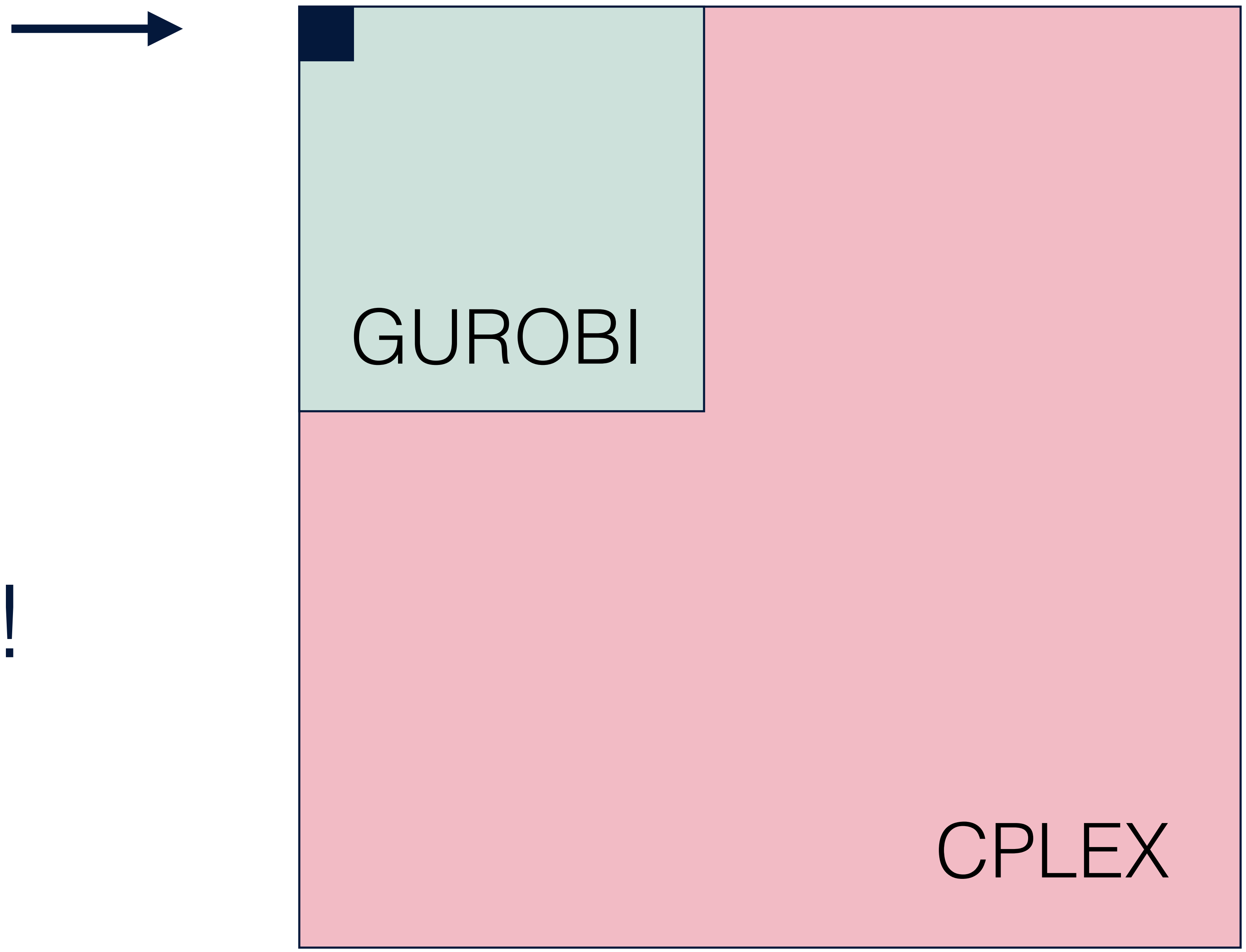
Library free

Detects Infeasibility

Embeddable

Compiled code size ~80kb

OSQP



300x
Reduction!

Interfaces

Languages

julia

python™

MATLAB®

Parsers

JuMP

CVXPY

YALMIP

OSQP interface

python

```python
# Create OSQP object
m = osqp.OSQP()

# Initialize solver
m.setup(P, q, A, l, u,
        settings)

# Solve
results = m.solve()

# Update cost with q_new
m.update(q=q_new)

# Solve again
results_new = m.solve()
```

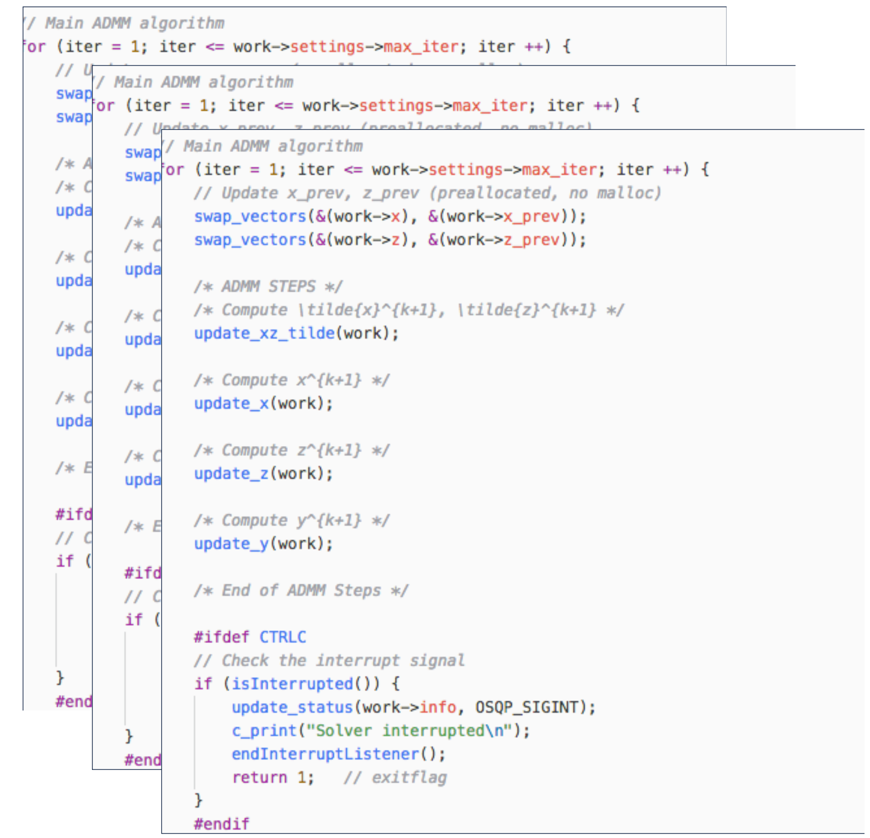
MATLAB

```matlab
% Create OSQP object
m = osqp();

% Initialize solver
m.setup(P, q, A, l, u,
        settings);

% Solve
results = m.solve();

% Update cost with q_new
m.update('q', q_new);

% Solve again
results_new = m.solve();
```

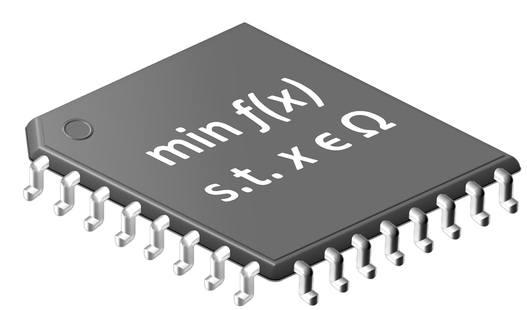
Code generation

```
# Create OSQP object
m = osqp.OSQP()

# Initialize solver
m.setup(P, q, A, l, u,
        settings)

# Generate C code
m.codegen('folder_name')
```

Optimized C code



Embedded hardware



Numerical Example

Lasso

minimize $\quad \|Ax - b\|_2^2 + \lambda \|x\|_1$

Features
n

Data points
$m = 100n$

Lasso

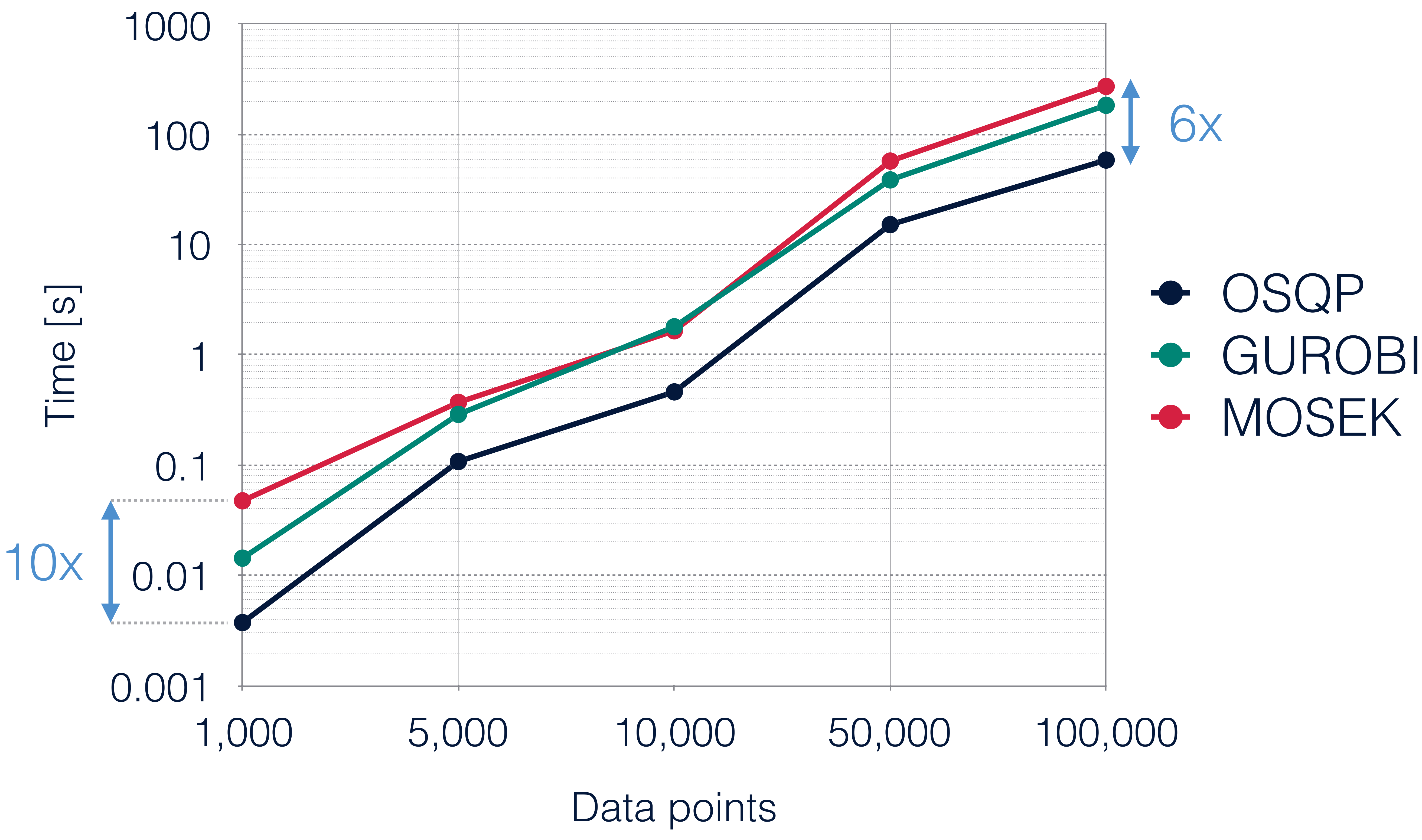
Weighting parameter

minimize $\quad \|Ax - b\|_2^2 + \lambda \|x\|_1$

Features
n

Data points
$m = 100n$

Lasso timings



OSQP
in
“meta-algorithms”

MIQP

$$
\begin{array}{ll}
\text{minimize} & \frac{1}{2}x^T P x + q^T x \\
\text{subject to} & l \leq Ax \leq u \\
 & x_i \in \mathbf{Z} \quad \forall i \in \mathbf{I}
\end{array}
$$

MIQP

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & l \leq Ax \leq u \\ & x_i \in \mathbf{Z} \quad \forall i \in \mathbf{I} \end{array}$$

Integer constraints

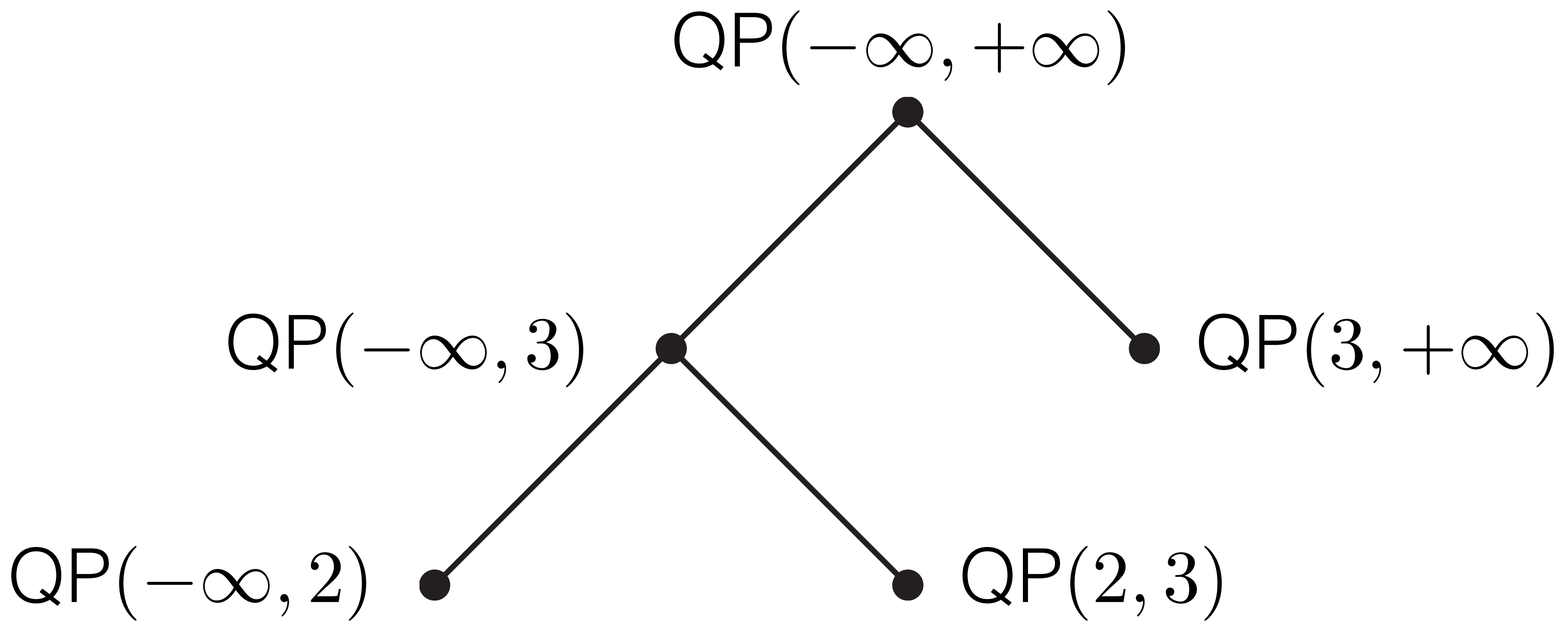
MIQP

minimize $\quad \frac{1}{2}x^TPx + q^Tx$

subject to $\quad l \leq Ax \leq u$

$\quad x_i \in \mathbf{Z} \quad \forall i \in \mathbf{I}$

Integer constraints



Inner QPs

$\text{QP}(\underline{x}, \overline{x})$

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & l \leq Ax \leq u \\ & \underline{x}_i \leq x_i \leq \overline{x}_i \quad \forall i \in \mathbf{I} \end{array}$$

Inner QPs

$\text{QP}(\underline{x}, \overline{x})$

$$
\begin{array}{ll}
\text{minimize} & \frac{1}{2} x^T P x + q^T x \\
\text{subject to} & l \leq Ax \leq u \\
& \underline{x}_i \leq x_i \leq \overline{x}_i \quad \forall i \in \mathbf{I}
\end{array}
$$

Changing bounds

↓

Reusing factorization

Saving computations

Factorization caching

+

Warm starting

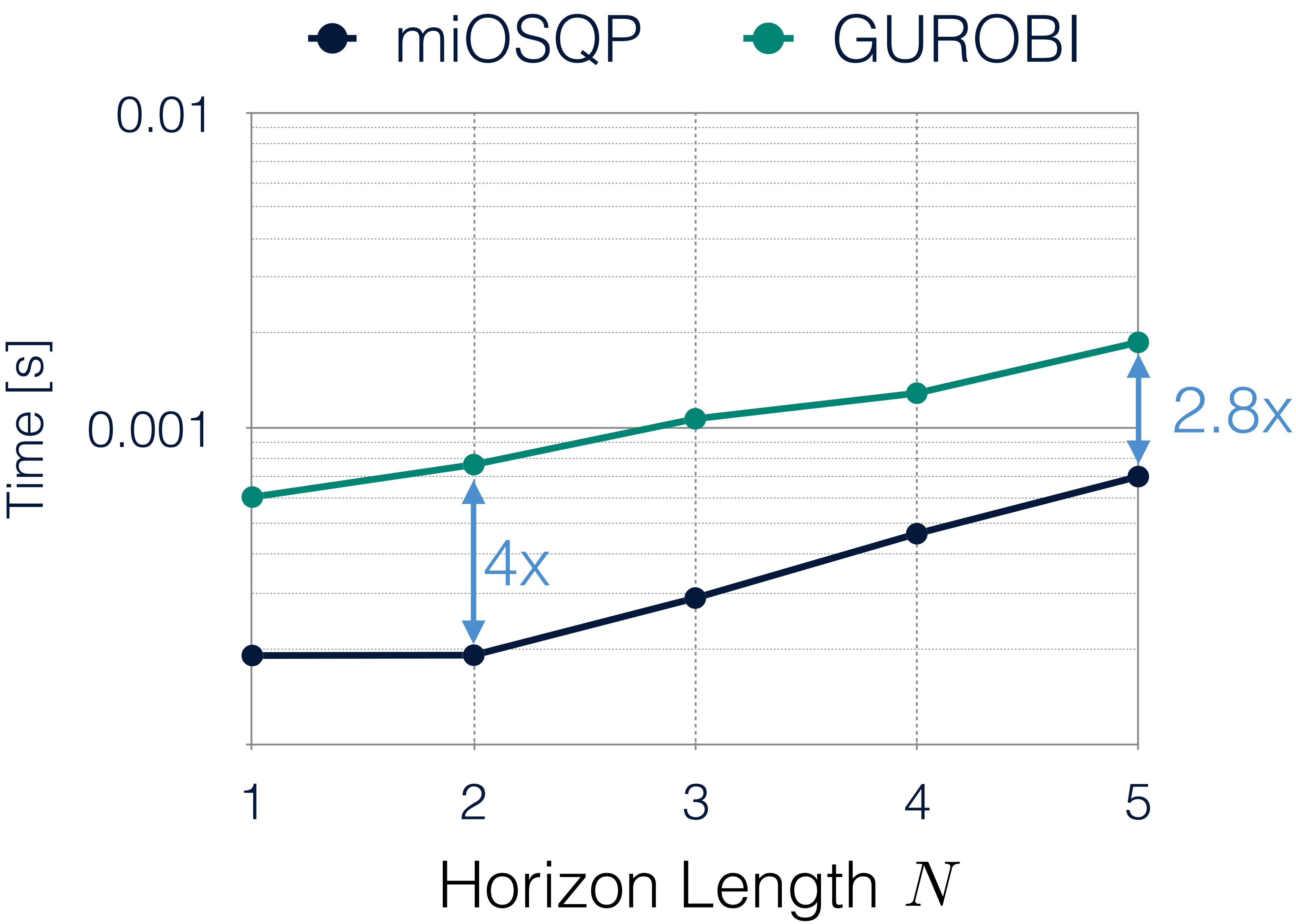
ADMM

QP$(\underline{x}, \overline{x})$

Repeated MIQPs

MIQP Timings

Hybrid Model Predictive Control



Simple Python Implementation!

Conclusions

Acknowledgements

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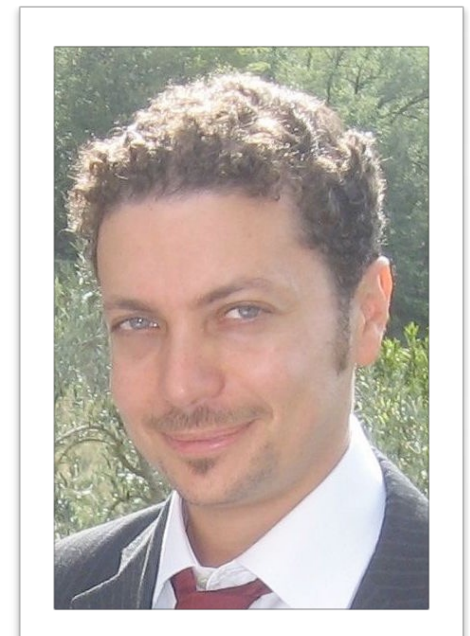
Oxford

Nicholas Moehle

Stanford

Paul Goulart

Oxford



Alberto Bemporad

IMT Lucca

Stephen Boyd

Stanford

Final remarks

OSQP

- Simple
- Robust
- Embeddable

Julia interface

- Exploit Initialization
- C code generation
- New high-level algorithms

References

B. Stellato, G. Banjac, P. Goulart, A. Bemporad and S. Boyd. ***OSQP: An Operator Splitting Solver for Quadratic Programs.*** *(Coming soon!)*

G. Banjac, P. Goulart, B. Stellato, and S. Boyd. ***Infeasibility detection in the alternating direction method of multipliers for convex optimization.*** optimization-online.org, 2017

G. Banjac, B.Stellato, N. Moehle, P. Goulart, A. Bemporad and S. Boyd. ***Embedded code generation using the OSQP solver.*** IEEE Conference on Decision and Control (CDC) (submitted), 2017

B. Stellato, V. Naik, A. Bemporad, P. Goulart, and S. Boyd. ***Embedded mixed-integer quadratic optimization using the OSQP solver.*** IEEE Conference on Decision and Control (CDC) (submitted), 2017